

Methods of proof

Contents

- 1 Direct proof
- 2 Disjunction of cases
- 3 Contraposition
- 4 Contradiction
- 5 By giving a counter example
- 6 Successive equivalences
- 7 Induction.

Introduction

- In mathematics, proving statements is fundamental.
- Different methods are used to establish the truth of mathematical statements.
- Here are some common methods of proofs :

Direct proof

- A direct proof demonstrates the truth of a proposition by a straightforward chain of logical deductions.
- We want to show that the proposition « $P \implies Q$ » is true.
- We assume that P is true and we show that then Q is true.
- This is the method you are most familiar with.

Example 1

- Let n be a natural number.
- Show that : n is even $\implies n^2$ is even.
- n is even $\implies \exists k \in \mathbb{N}$ such that $n = 2k$
- We have $n^2 = (2k)^2 = 2(2k^2) = 2l$ with $l = 2k^2 \in \mathbb{N}$.
- And consequently n^2 is even.

Example 2

- Show that : $x, y \in]-1, 1[\implies \frac{x+y}{1+xy} \in]-1, 1[.$

- We have

$$\alpha \in]-1, 1[\iff -1 < \alpha < 1 \iff |\alpha| < 1 \iff \alpha^2 < 1 \iff \alpha^2 - 1 < 0$$

- Let $x, y \in]-1, 1[.$

$$\begin{aligned} \left(\frac{x+y}{1+xy} \right)^2 - 1 &= \frac{(x+y)^2 - (1+xy)^2}{(1+xy)^2} = \frac{x^2 + y^2 - 1 - x^2y^2}{(1+xy)^2} \\ &= \frac{x^2 - 1 + y^2(1-x^2)}{(1+xy)^2} = \frac{(x^2-1)(1-y^2)}{(1+xy)^2} < 0 \end{aligned}$$

- because $x, y \in]-1, 1[$, and $x^2 < 1, y^2 < 1$

- then $\left(\frac{x+y}{1+xy} \right)^2 - 1 < 0 \implies \left(\frac{x+y}{1+xy} \right)^2 < 1 \implies \left| \frac{x+y}{1+xy} \right| < 1$

- Finally : $\left(\frac{x+y}{1+xy} \right)^2 - 1 < 0 \implies \frac{x+y}{1+xy} \in]-1, 1[.$

Disjunction of cases

- If we want to check a proposition $P(x)$ for all x in a set E ,
- we show the proposition for the x in a part A of E ,
- then for the x not belonging to A .
- This is the method of disjunction of cases or case by case.

Example 1

- Let $n \in \mathbb{N}$. Show that : $n(n+1)(n+2)$ is even.

First case : n is even $\exists k \in \mathbb{N}$ such that $n = 2k$

$$n(n+1)(n+2) = 2k(2k+1)(2k+2) = 2l \text{ with} \\ l = k(2k+1)(2k+2) \in \mathbb{N}$$

- therefore $n(n+1)(n+2)$ is even.

Second case : n is odd $\exists k \in \mathbb{N}$ such that $n = 2k + 1$

$$n(n+1)(n+2) = (2k+1)(2k+2)(2k+3) \\ = 2(2k+1)(k+1)(2k+3) \\ = 2l \text{ with } l = (2k+1)(k+1)(2k+3) \in \mathbb{N}.$$

- So, $n(n+1)(n+2)$ is even.
- Conclusion : $\forall n \in \mathbb{N}$, $n(n+1)(n+2)$ is even.

Example 2

- Show that :

$$\forall x \in \mathbb{R}, |x - 2| \leq x^2 - 3x + 3$$

- We have $|x - 2| = \begin{cases} x - 2 & \text{if } x \geq 2 \\ 2 - x & \text{if } x < 2 \end{cases}$



$$\text{If } x \geq 2, |x - 2| = x - 2 \leq x^2 - 3x + 3$$



then $x^2 - 4x + 5 \geq 0$ it's true because $\Delta = -4 < 0$



$$\text{If } x < 2, |x - 2| = 2 - x \leq x^2 - 3x + 3$$



then $x^2 - 2x + 1 = (x - 1)^2 \geq 0$ it is true.

- Conclusion :

$$\forall x \in \mathbb{R}, |x - 2| \leq x^2 - 3x + 3$$

Contraposition

- A proof by contrapositive proves an implication $P \implies Q$ by proving the equivalent contrapositive statement $\overline{Q} \implies \overline{P}$.
- Proof by contraposition is based on the following equivalence

$$(P \implies Q) \iff (\overline{Q} \implies \overline{P})$$

- So, if we want to show the proposition $P \implies Q$
- We actually show that if \overline{Q} is true then \overline{P} is true.

Example 1

- Let $n \in \mathbb{N}$.
- Show that : n^2 is even $\implies n$ is even.
- n is odd $\implies \exists k \in \mathbb{N}$ such that $n = 2k + 1$;
- We have

$$n^2 = (2k + 1)^2 = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1 = 2l + 1$$

- with $l = 2k^2 + 2k \in \mathbb{N}$.
- And consequently n^2 is odd .
- Conclusion : n^2 is even $\implies n$ is even.

Example 2

- Let $x, y \in \mathbb{R}$.
- Show that :

$$x \neq y \text{ and } xy \neq 1 \implies \frac{x}{x^2 + x + 1} \neq \frac{y}{y^2 + y + 1}$$

Example 2

- $$\frac{x}{x^2 + x + 1} = \frac{y}{y^2 + y + 1} \implies x(y^2 + y + 1) = y(x^2 + x + 1)$$

Contradiction

- A proof by contradiction assumes that the statement to be proven is false and then shows that this assumption leads to a contradiction.
- Let R be a proposition. We know that $R \vee \bar{R}$ is true.
- To show that R is true, we assume that R is false, that is to say \bar{R} is true and we show that we obtain a contradiction.
- If R is an implication, $R \approx P \implies Q$
- We have $\overline{P \implies Q} \iff P \wedge \bar{Q}$
- Proof by contradiction to show that $P \implies Q$ is based on the following principle :
- we assume both that P is true and that Q is false and we look for a contradiction.
- So if P is true then Q must be true and therefore $P \implies Q$ is true.

Example 1

- Let show by contradiction that $\sqrt{2}$ is irrational (not rational) :
- We assume that $\sqrt{2} \in \mathbb{Q}$.
- $\sqrt{2} \in \mathbb{Q} \implies \sqrt{2} = \frac{a}{b}$ with a and b are natural numbers that are prime to each other. (the fraction $\frac{a}{b}$ is irreducible).

\implies

Example 2

- Let $a, b > 0$.
- Show that :

$$\frac{a}{1+b} = \frac{b}{1+a} \implies a = b$$

Example 2

- Assume that $\frac{a}{1+b} = \frac{b}{1+a}$ with $a \neq b$
- $\frac{a}{1+b} = \frac{b}{1+a} \implies a(1+b) = b(1+a)$
 \implies